#### Introduction to **Parallel Programming & Cluster Computing MPI Collective Communications**

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## Q

#### **Point to Point Always Works**

- MPI\_Send and MPI\_Recv are known as <u>*Point to Point*</u> communications: they communicate from one MPI process to another MPI process.
- But, what if you want to communicate like one of these?
  - one to many
  - many to one
  - many to many
- These are known as *collective communications*.
- MPI\_Send and MPI\_Recv can accomplish any and all of these but should you use them that way?



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#### Point to Point Isn't Always Good

- We're interested in *collective communications*:
  - one to many
  - many to one
  - many to many
- In principle, MPI\_Send and MPI\_Recv can accomplish any and all of these.
- But that may be:
  - inefficient;
  - inconvenient and cumbersome to code.
- So, the designers of MPI came up with routines that perform these collective communications for you.





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#### **Collective Communications**

- MPI\_Bcast
- MPI\_Reduce, MPI\_Allreduce
- MPI\_Gather, MPI\_Gatherv, MPI\_Allgather, MPI\_Allgatherv
- MPI\_Scatter, MPI\_Scatterv
- MPI\_Alltoall, MPI\_Alltoallv









What happens if one process has data that everyone else needs to know?

For example, what if the server process needs to send a value that it input from standard input to the other processes?

MPI\_Bcast(&length, 1, MPI\_INTEGER,

source, MPI COMM WORLD);

Notice:

- MPI\_Bcast doesn't use a tag.
- The call is the same for both the sender and all of the receivers (<u>COUNTERINTUITIVE!</u>).

All processes have to call **MPI\_Bcast** at the same time; everyone waits until everyone is done.







#### **Broadcast Example Part 1**

```
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <mpi.h>
int main (int argc, char** argv)
{ /* main */
  const int server = 0;
  const int source = server;
  float* array = (float*)NULL;
  int
         length, index;
         number of processes, my rank, mpi error code;
  int
  mpi error code = MPI Init(&argc, &argv);
  mpi error code = MPI Comm rank(MPI COMM WORLD, &my rank);
  mpi error code = MPI Comm size(MPI COMM WORLD,
                                  &number of processes);
```





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#### **Broadcast Example Part 2**

```
if (my rank == source) {
   scanf("%d", &length);
  } /* if (my rank == source) */
 fprintf(stderr, "%d: before MPI Bcast, length = %d\n",
   my rank, length);
 mpi error code =
   MPI Bcast(&length, 1, MPI INTEGER, source, MPI COMM WORLD);
 fprintf(stderr, "%d: after MPI Bcast, length = %d\n",
   my rank, length);
 array = (float*)malloc(sizeof(float) * length);
 if (my rank == source) {
   for (index = 0; index < length; index++) {</pre>
     array[index] = sqrt(index * 1.0); /* Or whatever you want */
   } /* for index */
 } /* if (my rank == source) */
 mpi error code =
   MPI Bcast(array, length, MPI FLOAT, source, MPI COMM WORLD);
 mpi error code = MPI Finalize();
} /* main */
```







#### **Broadcast Compile & Run**

<pre>% mpicc -o mpibroadcast mpibroadcast.c -lm</pre>				
<pre>% mpirun -np 8 mpibroadcast</pre>				
4:	before	MPI_Bcast,	length =	0
7:	before	MPI_Bcast,	length =	0
3:	before	MPI_Bcast,	length =	0
5:	before	MPI_Bcast,	length =	0
6:	before	MPI_Bcast,	length =	0
2:	before	MPI_Bcast,	length =	0
0:	before	MPI_Bcast,	length =	1000000
0:	after	MPI_Bcast,	length =	1000000
2:	after	MPI_Bcast,	length =	1000000
4:	after	MPI_Bcast,	length =	1000000
5:	after	MPI_Bcast,	length =	1000000
7:	after	MPI_Bcast,	length =	1000000
6:	after	MPI_Bcast,	length =	1000000
3:	after	MPI_Bcast,	length =	1000000
1:	before	MPI_Bcast,	length =	0
1:	after	MPI Bcast,	length =	1000000









#### Reductions

A <u>*reduction*</u> converts an array to a scalar: for example, sum, product, minimum value, maximum value, Boolean AND, Boolean OR, etc.

- Reductions are so common, and so important, that MPI has two routines to handle them:
- **MPI\_Reduce**: sends result to a single specified process
- MPI\_Allreduce: sends result to all processes (and therefore may take longer)







#### **Reduction Example Part 1**

```
#include <stdio.h>
#include <stdlib.h>
#include <mpi.h>
int main (int argc, char** argv)
{ /* main */
 const int server
                        = 0;
  const int destination = server;
  int value, value sum;
  int number of processes, my rank, mpi error code;
 mpi error code = MPI Init(&argc, &argv);
 mpi error code = MPI Comm rank(MPI COMM WORLD, &my rank);
 mpi error code = MPI Comm size (MPI COMM WORLD,
                                  &number of processes);
```







#### **Reduction Example Part 1**

```
value = my rank * number of processes;
  fprintf(stderr, "%d: reduce
                                value
                                          = %d n'',
   my rank, value);
 mpi error code =
   MPI Reduce (&value, &value sum, 1, MPI INT, MPI SUM,
                 destination, MPI COMM WORLD);
 fprintf(stderr, "%d: reduce value sum = %d\n",
   my rank, value sum);
 mpi error code =
   MPI Allreduce(&value, &value sum, 1, MPI INT, MPI SUM,
                              MPI COMM WORLD);
 fprintf(stderr, "%d: allreduce value sum = %d\n",
   my rank, value sum);
 mpi error code = MPI Finalize();
} /* main */
```







#### **Reduce: Compiling and Running**

- % mpicc -o mpireduce mpireduce.c
- % mpirun -np 8 mpireduce
- 0: reduce value = 0 4: reduce value = 32
- 6: reduce value = 48
- 7: reduce value = 56
- 3: reduce value = 24
- 2: reduce value = 16 5: reduce value = 40
- 1: reduce value = 8
- 7: reduce value\_sum = -9120
- 3: reduce value\_sum = -9120 2: reduce value sum = -9120
- 5: reduce value sum = -9120
- 1: reduce value sum = -9120
- 6: reduce value sum = -9120
- 4: reduce value sum = -9120
- 0: reduce value sum = 224

- 2: allreduce value\_sum = 224
- 7: allreduce value\_sum = 224
- 4: allreduce value\_sum = 224
- 3: allreduce value\_sum = 224
- 1: allreduce value\_sum = 224
- 5: allreduce value\_sum = 224
- 0: allreduce value\_sum = 224
- 6: allreduce value\_sum = 224







### Why Two Reduction Routines?

MPI has two reduction routines because of the high cost of each communication.

- If only one process needs the result, then it doesn't make sense to pay the cost of sending the result to all processes.
- But if all processes need the result, then it may be cheaper to reduce to all processes than to reduce to a single process and then broadcast to all.
- You can think of MPI\_Allreduce as MPI\_Reduce followed by MPI\_Bcast (though it doesn't have to be implemented that way).







#### **Reduction on Arrays #1**

MPI\_Reduce and MPI\_Allreduce are actually designed to work on arrays, where the corresponding elements of each source array are reduced into the corresponding element of the destination array (all of the same length):

For example:







#### **Reduction on Arrays #2**

global\_force\_on\_particle[p] =
 local\_force\_on\_particle[p] on Rank 0 +
 local\_force\_on\_particle[p] on Rank 1 +
 local\_force\_on\_particle[p] on Rank 2 +

• • •

local\_force\_on\_particle[p] on Rank np-1;







#### **Scatter and Gather**

- To *scatter* is to send data from one place to many places.
- To *gather* is to receive data from many places into one place.
- MPI has a variety of scatter and gather routines:
  - MPI\_Scatter, MPI\_Scatterv
  - MPI\_Gather, MPI\_Gatherv, MPI\_Allgather, MPI\_Allgatherv
- The scatter routines split up a single larger array into smaller subarrays, one per MPI process, and send each subarray to an MPI process.
- The gather routines receive many smaller subarrays, one per MPI process, and assemble them into a single larger array.











MPI\_Scatter takes an array whose length is divisible by the number of MPI processes, and splits it up into subarrays of equal length, then sends one subarray to each MPI process.

So, for a large array of length 100 on 5 MPI processes:

- each smaller subarray has length 20;
- large\_array[0] .. large\_array[19] go to small\_array on Rank 0;
- large\_array[20]..large\_array[39] go to small\_array on Rank 1;

etc









MPI\_Scatterv is just like MPI\_Scatter, except that the subarray lengths don't have to be the same (and therefore the length of the large array doesn't have to be divisible by the number of MPI processes).

The displacements array says where each small subarray begins within the large array.









large\_array, large\_array\_length,
MPI\_DATATYPE, destination, MPI\_COMMUNICATOR);

So, for a small subarray of length 20 on each of 5 MPI processes:

- the large array on the destination process has length 100;
- large\_array[0] .. large\_array[19] come from small\_array on Rank 0;
- large\_array[20]..large\_array[39] come from small\_array on Rank 1;









MPI\_Gatherv is just like MPI\_Gather, except that the subarray lengths don't have to be the same (and therefore the length of the large array doesn't have to be divisible by the number of MPI processes).

MPI\_Gatherv(small\_subarray, small\_subarray\_length,

MPI\_DATATYPE,

large\_array, small\_subarray\_lengths,

displacements,

MPI\_DATATYPE, destination, MPI\_COMMUNICATOR);

The displacements array says where each small subarray begins within the large array.





#### MPI\_Allgather & MPI\_Allgatherv

MPI\_Allgather and MPI\_Allgatherv are the same as MPI\_Gather and MPI\_Gatherv, except that the large array gets filled on every MPI process, so no destination process argument is needed.





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