Introduction to Parallel Programming & Cluster Computing

MPI Collective Communications

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Point to Point Always Works

- MPI_Send and MPI_Recv are known as **Point to Point** communications: they communicate from one MPI process to another MPI process.
- But, what if you want to communicate like one of these?
  - one to many
  - many to one
  - many to many
- These are known as **collective communications**.
- MPI_Send and MPI_Recv can accomplish any and all of these – but should you use them that way?
We’re interested in collective communications:
- one to many
- many to one
- many to many

In principle, MPI_Send and MPI_Recv can accomplish any and all of these.

But that may be:
- inefficient;
- inconvenient and cumbersome to code.

So, the designers of MPI came up with routines that perform these collective communications for you.
Collective Communications

- MPI_Bcast
- MPI_Reduce, MPI_Allreduce
- MPI_Gather, MPI_Gatherv, MPI_Allgather, MPI_Allgatherv
- MPI_Scatter, MPI_Scatterv
- MPI_Alltoall, MPI_Alltoallv
What happens if one process has data that everyone else needs to know?

For example, what if the server process needs to send a value that it input from standard input to the other processes?

```c
MPI_Bcast(&length, 1, MPI_INTEGER, source, MPI_COMM_WORLD);
```

Notice:

- **MPI_Bcast** doesn’t use a tag.
- The call is the same for both the sender and all of the receivers (**COUNTERINTUITIVE**!).

All processes have to call **MPI_Bcast** at the same time; everyone waits until everyone is done.
Broadcast Example Part 1

```c
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <mpi.h>

int main (int argc, char** argv)
{ /* main */
    const int server = 0;
    const int source = server;
    float* array   = (float*)NULL;
    int   length, index,
    int   number_of_processes, my_rank, mpi_error_code;

    mpi_error_code = MPI_Init(&argc, &argv);
    mpi_error_code = MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);
    mpi_error_code = MPI_Comm_size(MPI_COMM_WORLD,
                                    &number_of_processes);
}
```
if (my_rank == source) {
    scanf("%d", &length);
} /* if (my_rank == source) */
fprintf(stderr, "%d: before MPI_Bcast, length = %d\n",
    my_rank, length);
mpi_error_code =
    MPI_Bcast(&length, 1, MPI_INTEGER, source, MPI_COMM_WORLD);
fprintf(stderr, "%d: after MPI_Bcast, length = %d\n",
    my_rank, length);
array = (float*)malloc(sizeof(float) * length);
if (my_rank == source) {
    for (index = 0; index < length; index++) {
        array[index] = sqrt(index * 1.0); /* Or whatever you want */
    } /* for index */
} /* if (my_rank == source) */
mpi_error_code =
    MPI_Bcast(array, length, MPI_FLOAT, source, MPI_COMM_WORLD);
mpi_error_code = MPI_Finalize();
} /* main */
Broadcast Compile & Run

% mpicc -o mpibroadcast mpibroadcast.c -lm
% mpirun -np 8 mpibroadcast

4: before MPI_Bcast, length = 0
7: before MPI_Bcast, length = 0
3: before MPI_Bcast, length = 0
5: before MPI_Bcast, length = 0
6: before MPI_Bcast, length = 0
2: before MPI_Bcast, length = 0
0: before MPI_Bcast, length = 1000000
0: after MPI_Bcast, length = 1000000
2: after MPI_Bcast, length = 1000000
4: after MPI_Bcast, length = 1000000
5: after MPI_Bcast, length = 1000000
7: after MPI_Bcast, length = 1000000
6: after MPI_Bcast, length = 1000000
3: after MPI_Bcast, length = 1000000
1: before MPI_Bcast, length = 0
1: after MPI_Bcast, length = 1000000
Reductions

A **reduction** converts an array to a scalar: for example, sum, product, minimum value, maximum value, Boolean AND, Boolean OR, etc.

Reductions are so common, and so important, that MPI has two routines to handle them:

- **MPI_Reduce**: sends result to a single specified process
- **MPI_Allreduce**: sends result to all processes (and therefore may take longer)
Reduction Example Part 1

```c
#include <stdio.h>
#include <stdlib.h>
#include <mpi.h>

int main (int argc, char** argv)
{
    /* main */
    const int server = 0;
    const int destination = server;
    int value, value_sum;
    int number_of_processes, my_rank, mpi_error_code;

    mpi_error_code = MPI_Init(&argc, &argv);
    mpi_error_code = MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);
    mpi_error_code = MPI_Comm_size(MPI_COMM_WORLD,
                                    &number_of_processes);
}
```
value = my_rank * number_of_processes;
fprintf(stderr, "%d: reduce value = %d\n", 
    my_rank, value);
mpi_error_code =
    MPI_Reduce (&value, &value_sum, 1, MPI_INT, MPI_SUM, 
        destination, MPI_COMM_WORLD);
fprintf(stderr, "%d: reduce value_sum = %d\n", 
    my_rank, value_sum);
mpi_error_code =
    MPI_Allreduce(&value, &value_sum, 1, MPI_INT, MPI_SUM, 
        MPI_COMM_WORLD);
fprintf(stderr, "%d: allreduce value_sum = %d\n", 
    my_rank, value_sum);
mpi_error_code = MPI_Finalize();
} /* main */
Reduce: Compiling and Running

```bash
% mpicc -o mpireduce mpireduce.c
% mpirun -np 8 mpireduce

0: reduce  value     = 0
4: reduce  value     = 32
6: reduce  value     = 48
7: reduce  value     = 56
3: reduce  value     = 24
2: reduce  value     = 16
5: reduce  value     = 40
1: reduce  value     = 8
7: reduce  value_sum = -9120
3: reduce  value_sum = -9120
2: reduce  value_sum = -9120
5: reduce  value_sum = -9120
1: reduce  value_sum = -9120
6: reduce  value_sum = -9120
4: reduce  value_sum = -9120
0: reduce  value_sum = 224
```

2: allreduce value_sum = 224
7: allreduce value_sum = 224
4: allreduce value_sum = 224
3: allreduce value_sum = 224
1: allreduce value_sum = 224
5: allreduce value_sum = 224
0: allreduce value_sum = 224
6: allreduce value_sum = 224
Why Two Reduction Routines?

MPI has two reduction routines because of the high cost of each communication.

If only one process needs the result, then it doesn’t make sense to pay the cost of sending the result to all processes.

But if all processes need the result, then it may be cheaper to reduce to all processes than to reduce to a single process and then broadcast to all.

You can think of `MPI_Allreduce` as `MPI_Reduce` followed by `MPI_Bcast` (though it doesn’t have to be implemented that way).
Reduction on Arrays #1

MPI_Reduce and MPI_Allreduce are actually designed to work on arrays, where the corresponding elements of each source array are reduced into the corresponding element of the destination array (all of the same length):

\[
\text{MPI_Allreduce(} \text{source_array, destination_array, number_of_array_elements, MPI_DATATYPE, MPI_OPERATION, MPI_COMMUNICATOR);}
\]

For example:

\[
\text{MPI_Allreduce(} \text{local_force_on_particle, global_force_on_particle, number_of_particles, MPI_FLOAT, MPI_SUM, MPI_COMM_WORLD);}
\]
Reduction on Arrays #2

MPI_Allreduce(local_force_on_particle, global_force_on_particle, number_of_particles, MPI_FLOAT, MPI_SUM, MPI_COMM_WORLD);

\[
global\_force\_on\_particle[p] = 
local\_force\_on\_particle[p] \text{ on Rank 0 } + 
local\_force\_on\_particle[p] \text{ on Rank 1 } + 
local\_force\_on\_particle[p] \text{ on Rank 2 } + 
\ldots 
local\_force\_on\_particle[p] \text{ on Rank np-1};
\]
Scatter and Gather

- To scatter is to send data from one place to many places.
- To gather is to receive data from many places into one place.
- MPI has a variety of scatter and gather routines:
  - MPI_Scatter, MPI_Scatterv
  - MPI_Gather, MPI_Gatherv, MPI_Allgather, MPI_Allgatherv

- The scatter routines split up a single larger array into smaller subarrays, one per MPI process, and send each subarray to an MPI process.
- The gather routines receive many smaller subarrays, one per MPI process, and assemble them into a single larger array.
MPI_Scatter takes an array whose length is divisible by the number of MPI processes, and splits it up into subarrays of equal length, then sends one subarray to each MPI process.

MPI_Scatter(large_array, small_array_length, MPI_DATATYPE, small_subarray, small_subarray_length, MPI_DATATYPE, source, MPI_COMMUNICATOR);

So, for a large array of length 100 on 5 MPI processes:

- each smaller subarray has length 20;
- large_array[0] .. large_array[19] go to small_array on Rank 0;
- large_array[20]..large_array[39] go to small_array on Rank 1;
- etc
MPI_Scatterv

MPI_Scatterv is just like MPI_Scatter, except that the subarray lengths don’t have to be the same (and therefore the length of the large array doesn’t have to be divisible by the number of MPI processes).

MPI_Scatterv(large_array, small_subarray_lengths, displacements, MPI_DATATYPE, small_subarray, small_subarray_lengths, MPI_DATATYPE, source, MPI_COMMUNICATOR);

The displacements array says where each small subarray begins within the large array.
MPI_Gather receives a small array on each of the MPI processes, all subarrays of equal length, and joins them into a single large array on the destination MPI process.

MPI_Gather(small_subarray, small_subarray_length, MPI_DATATYPE, large_array, large_array_length, MPI_DATATYPE, destination, MPI_COMMUNICATOR);

So, for a small subarray of length 20 on each of 5 MPI processes:

- the large array on the destination process has length 100;
- large_array[0] .. large_array[19] come from small_array on Rank 0;
- large_array[20]..large_array[39] come from small_array on Rank 1;
- etc
MPI_Gatherv is just like MPI_Gather, except that the subarray lengths don’t have to be the same (and therefore the length of the large array doesn’t have to be divisible by the number of MPI processes).

MPI_Gatherv(small_subarray, small_subarray_length, MPI_DATATYPE, large_array, small_subarray_lengths, displacements, MPI_DATATYPE, destination, MPI_COMMUNICATOR);

The displacements array says where each small subarray begins within the large array.
MPI_Allgather & MPI_Allgatherv

MPI_Allgather and MPI_Allgatherv are the same as MPI_Gather and MPI_Gatherv, except that the large array gets filled on every MPI process, so no destination process argument is needed.
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